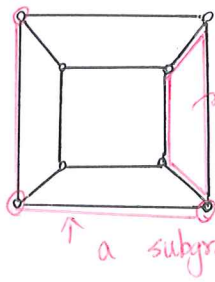


# Graph theory : trees and searching algorithms

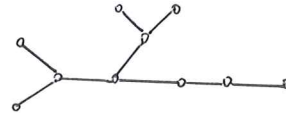


a cycle in a graph

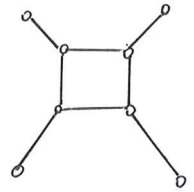
- A tree is a connected graph without cycle.

Property of a tree on  $n$  vertices :

- connected
- without cycle
- $n-1$  edges.



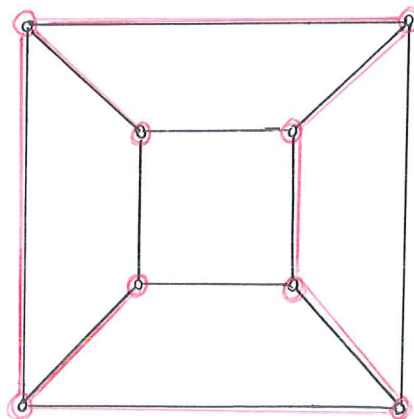
tree



not tree

- Any two properties out of the three defines a tree.
- A spanning tree of  $G$  is a subgraph of  $G$  using all vertices.
- A tree is the minimal connected structure.

in terms of  
number of edges



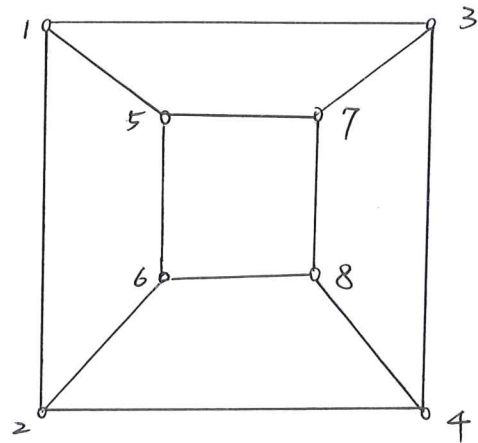
spanning tree

## Breadth first search

Search each vertex of a graph such that the algorithm searches all neighbors first before move to other vertices.

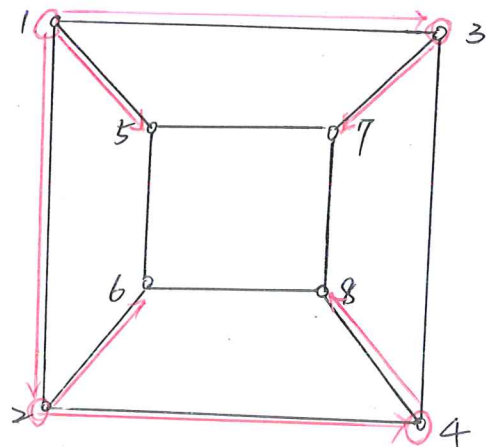
Exmp.

- Start from 1  
anchor at 1.
- Search 2, 3, 5,  
anchor at 2.
- Search 4, 6  
anchor at 3.
- Search 7,  
anchor at 5, 4
- Search 8,  
anchor at 6, 7, 8.



## Applications

- Search through vertices.
- Test connected or not.
- Generate a spanning tree.  
(ditree)



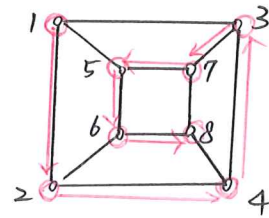
## Depth first search

Search each vertex of a graph such that the algorithm searches as deep as possible.

## Exmp

- Start from 1
  - 2 → 4 → 3
  - 7 → 5 → 6 → 8

Done!



Similar applications as BFS,  
same time complexity.

## Implementing the BFS

Pseudocode :

Input :

$G$ : graph,  $V$ : vertex

Output : an order by BFS.

searched = [v]

waiting = [v]

while waiting not empty:

anchor = first element in waiting

(and remove anchor from waiting)

new = [nbrs of anchor that is not yet searched]

search + = new

waiting + = new

Exmp (V=1)

Step	searched	waiting	anchor	new
0	[1]	[1]		
1	[1, 2, 3, 5]	[2, 3, 5]	1	[2, 3, 5]
2	[1, 2, 3, 5, 4, 6]	[3, 5, 4, 6]	2	[4, 6]
3	[1, 2, 3, 5, 4, 6, 7]	[5, 4, 6, 7]	3	[7]
4	same	[4, 6, 7]	5	[ ]
5	[1, 2, 3, 5, 4, 6, 7, 8]	[6, 7, 8]	4	[8]
6	done	[7, 8]	6	[ ]
7	done	[8]	7	[ ]
8	done	[ ]	8	[ ]

