Comparability and cocomparability bigraphs

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Dec 8, 2018
2018 TMS Annual Meeting, Taipei, Taiwan

Joint work with

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Colorado State U

(photos from department websites and personal websites)

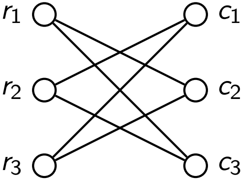
$\diagup$-free

$$\diagup = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

A $0,1$-matrix is $\diagup$-free if the rolws and columns can be permuted (independently) so that the resulting matrix does not contain $\diagup$ as a submatrix.

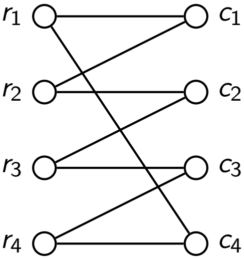
- How to recognize a $\diagup$-free matrix?
- If a matrix is $\diagup$-free, how to find the correct permutations?

C_6 is $\diagup$-free



$$\begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix} \xrightarrow{r_1 \leftrightarrow r_2} \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \end{bmatrix} \xrightarrow{c_2 \leftrightarrow c_3} \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{bmatrix} \xrightarrow{r_2 \leftrightarrow r_3} \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

C_8 is not $\diagup$-free



$$\begin{bmatrix} 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

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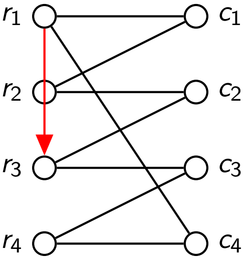

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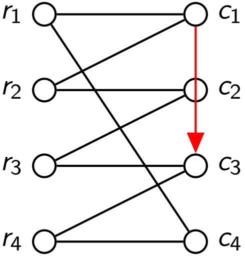

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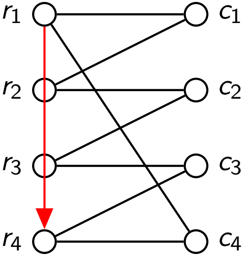

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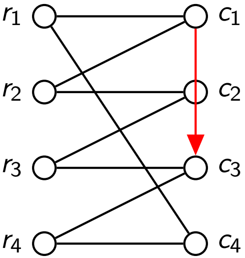

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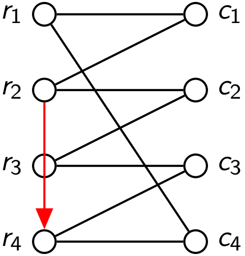

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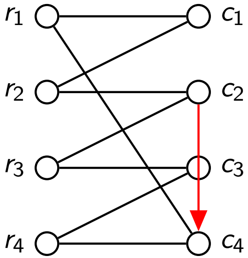

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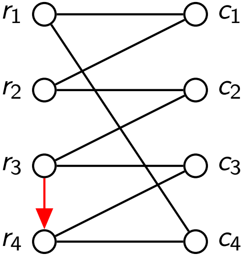

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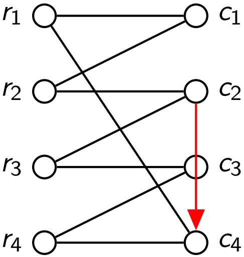

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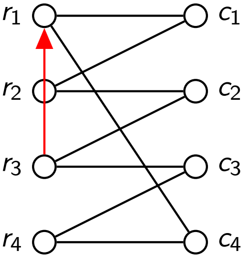


$$\begin{bmatrix} 1 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}$$

Submatrix avoiding problem

- A: a $0, 1$-matrix (usually the biadjacency matrix)
- S: a small matrix (usually 2×2)
- A is S-free if rows and columns can be permuted independently to avoid S.

How to recognize? How to realize?

- A: a square $0, 1$-matrix (usually the adjacency matrix)
- S: a small matrix (usually 2×2)
- A is symmetrically S-free if rows and columns can be permuted symmetrically to avoid S.

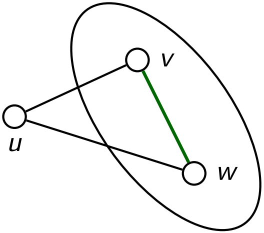
Chordal graph

- $R(G) = A(G) + I$, the neighborhood matrix
- principal $\Gamma = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$ with one of 1 embedded on the diagonal of the target matrix.

A graph G is called a chordal graph if one of the following holds.

- G has no long (≥ 4) cycle.
- $R(G)$ is symmetrically principal Γ-free.

$$\begin{array}{c|ccc} & u & v & w \\ \hline u & 1 & 1 & 1 \\ v & 1 & 1 & 1 \\ w & 1 & 1 & 1 \end{array}$$

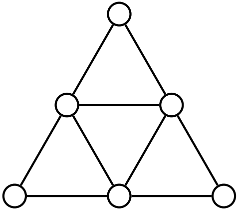


Strongly chordal graph

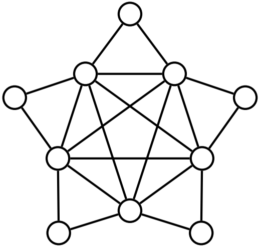
Theorem (Chang 1982; Farber 1983)

A graph G is called a strongly chordal graph if one of the following holds.

- *G has no long cycle and no induced trampoline.*
- *$R(G)$ is symmetrically ~~principal~~ Γ-free.*



trampolines



Chordal bigraph

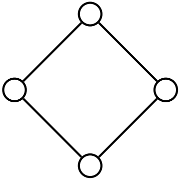
Theorem

A bipartite graph G is called a chordal bigraph if one of the following holds.

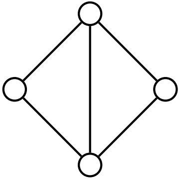
- *G has no long (≥ 6) cycle.*
- *The biadjacency matrix of G is ~~symmetrically~~ Γ-free.*

Comparability graph

A comparability graph is the graph of a poset, where two vertices are adjacent if and only if they are comparable.



poset



a comparability graph

On an induced P_3, the orientations of its two edges are related.

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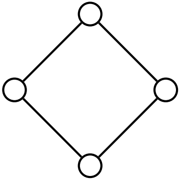

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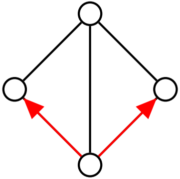

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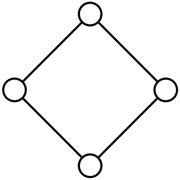

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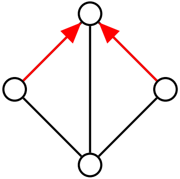

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On an induced P_3, the orientations of its two edges are related.

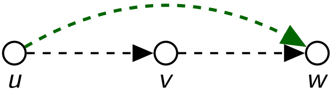
Cocomparability graph

- principal $\diagup = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$ with one of 1 embedded on the diagonal of the target matrix.

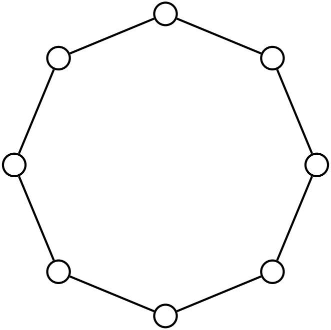
A graph G is called a cocomparability graph if one of the following holds.

- $\overline{G}$ is a comparability graph.
- $R(G) = A(G) + I$ is symmetrically principal $\diagup$-free.
- G has no invertible pair.
- G has no vertex asteroid. [Gilmore and Hoffman 1964]

$$\begin{array}{c|ccc} & u & v & w \\ \hline u & 1 & 0 & 0 \\ v & 0 & 1 & 0 \\ w & 1 & 0 & 1 \end{array}$$



Invertible pair: C_8 is not a cocomparability graph



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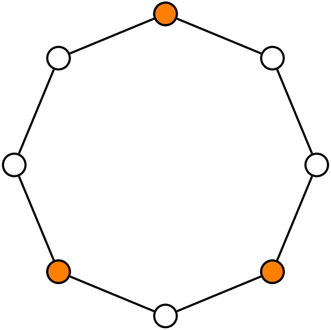

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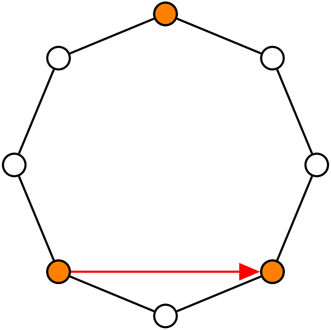

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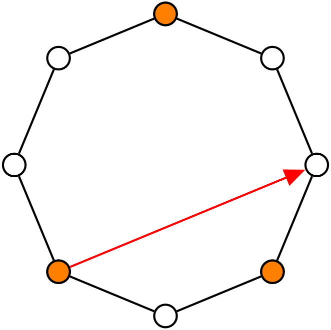

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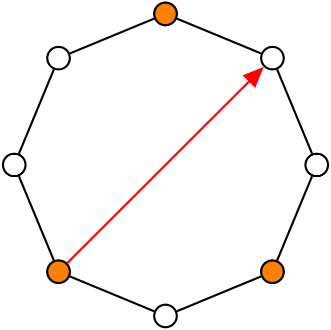

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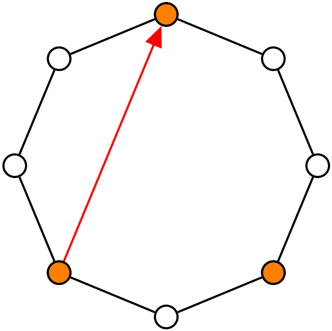

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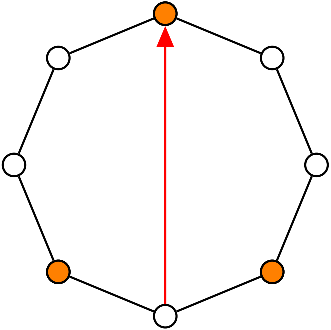

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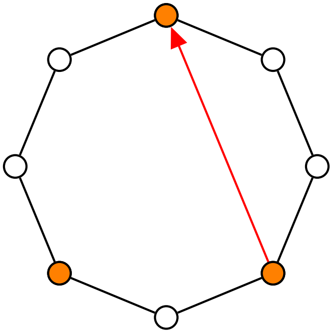

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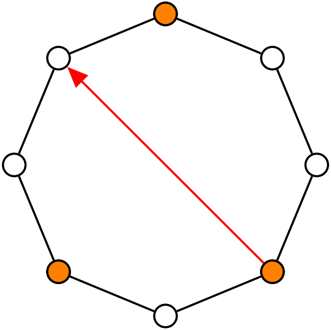

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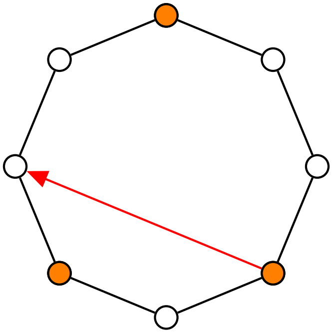

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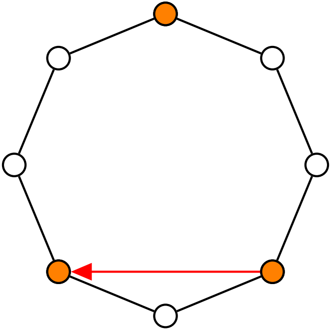

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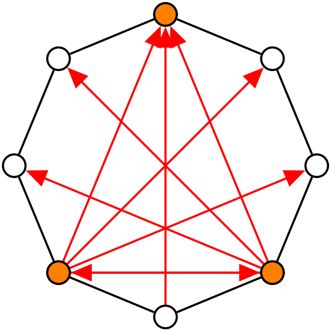
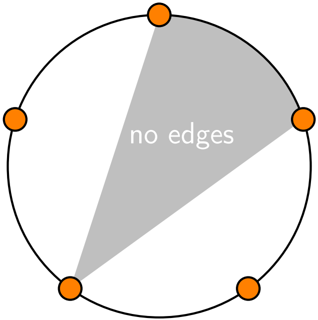


Vertex asteroid

Vertex asteroid: a set of vertices $v_0, \ldots, v_{2k}$ of odd numbers such that

- there is a path P_i from v_i to v_{i+1}, and
- v_i is not adjacent to any vertex on P_{i+k}.



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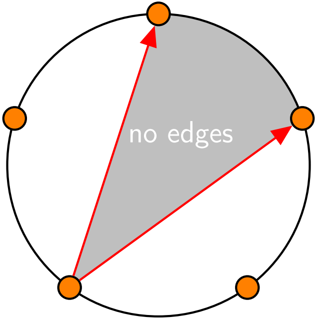

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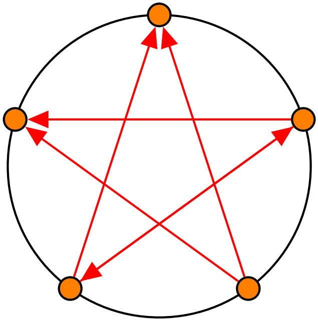
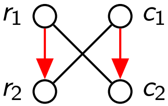


Cocomparability bigraph

A bipartite graph G is called a cocomparability bigraph if one of the following holds.

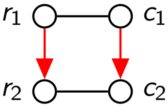
- the biadjacency matrix of G is ~~symmetrically principal~~ $\diagup$-free.
- G has no invertible pair.
- G has no edge asteroid.

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$



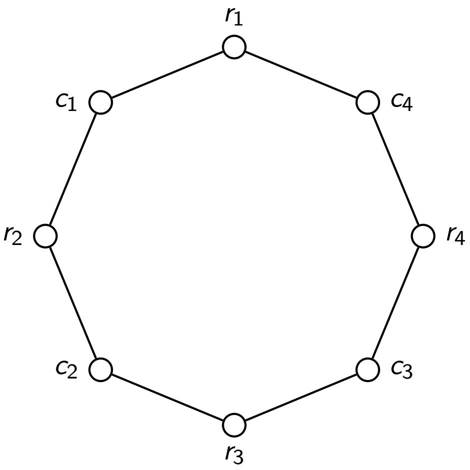
not allowed

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

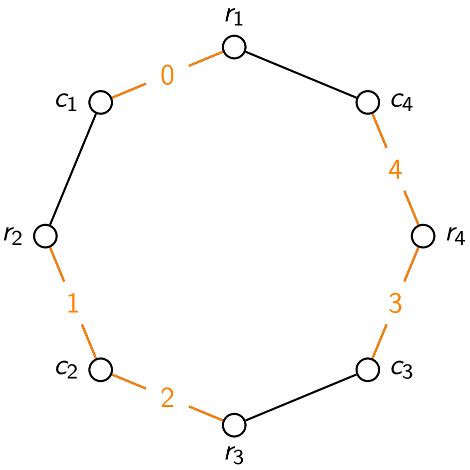


allowed

Invertible pair: C_8 is not $\diagup$-free (revisit)



Invertible pair: C_8 is not ⁄-free (revisit)



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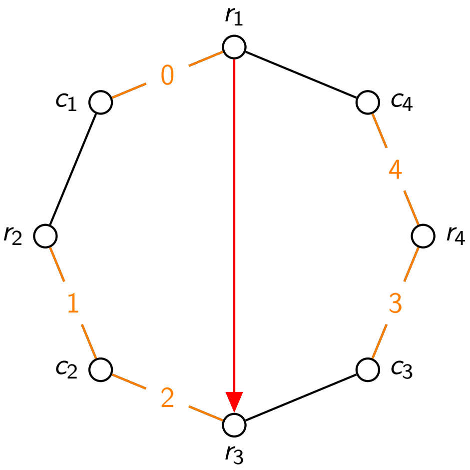

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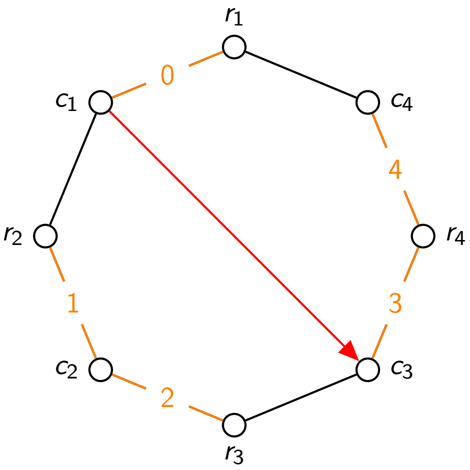

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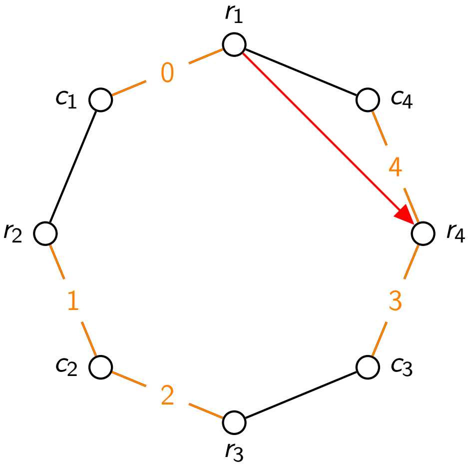

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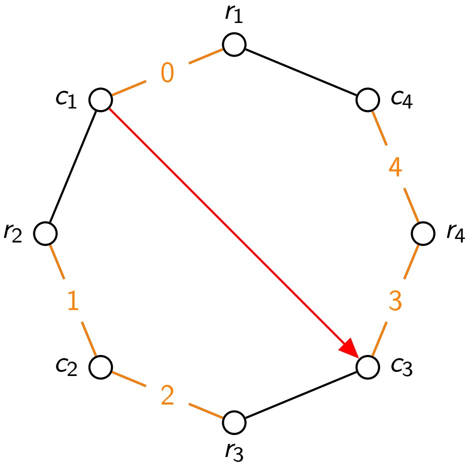

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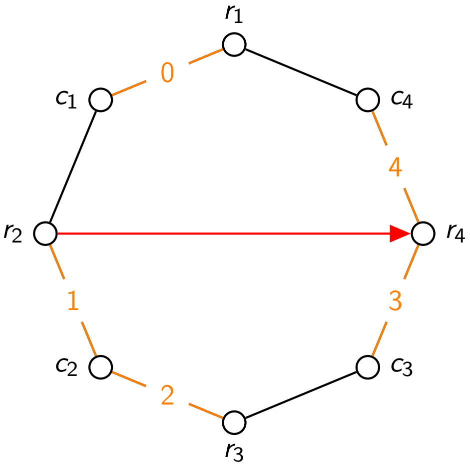

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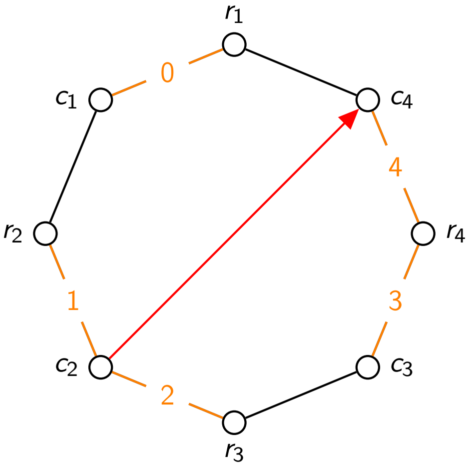

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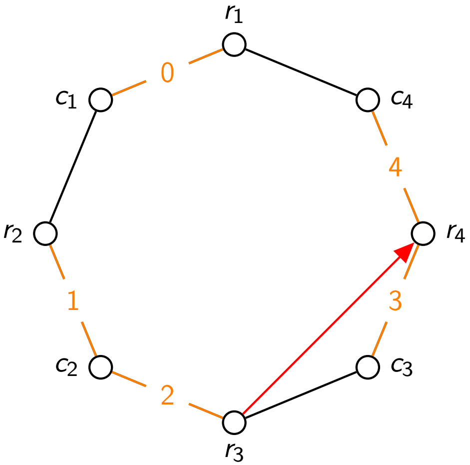

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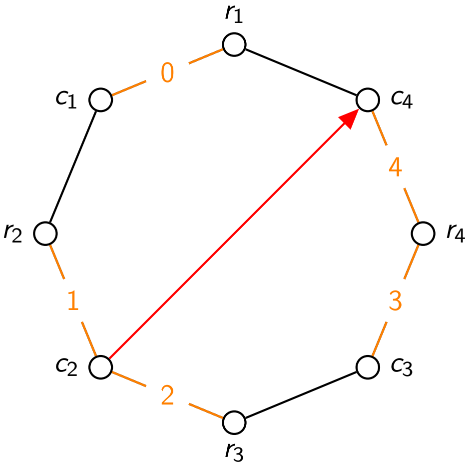

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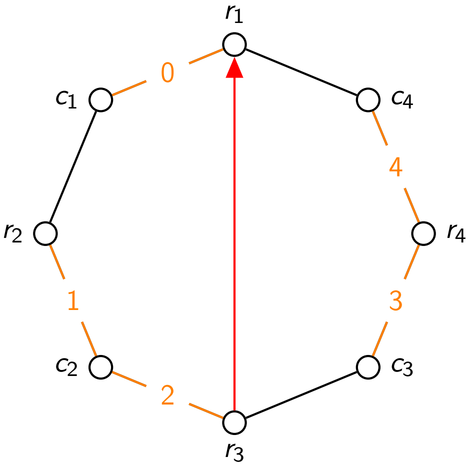

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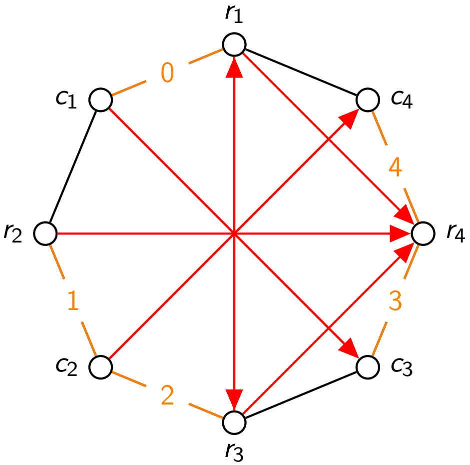
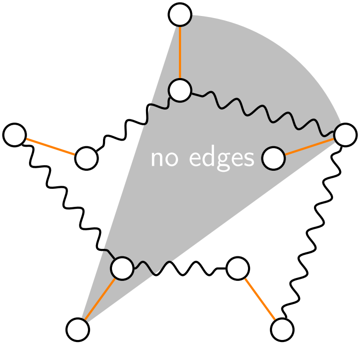


Edge asteroid

Edge asteroid: a set of edges $e_0, \ldots, e_{2k}$ of odd numbers such that

- there is a path P_i from e_i to e_{i+1}, and
- v_i is not adjacent to any vertex on P_{i+k} (including all vertices on e_{i+k} and e_{i+k+1}).



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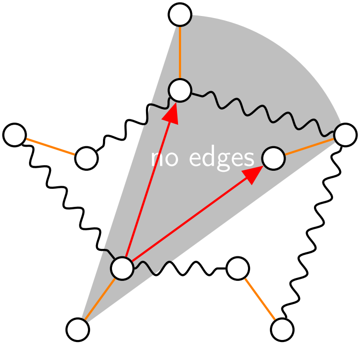

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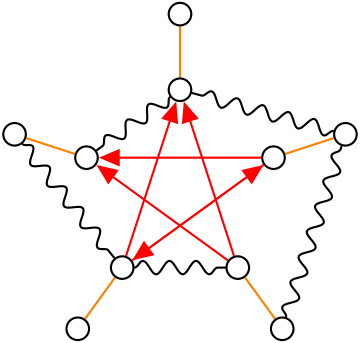
- v_i is not adjacent to any vertex on P_{i+k} (including all vertices on e_{i+k} and e_{i+k+1}).



Algorithms

Theorem (Hell, Huang, JL, and McConnell 2018+)

There is a polynomial-time algorithm to recognize a $\diagup$-free matrix.

Theorem (Hell, Huang, JL, and McConnell 2018+)

There is a polynomial-time algorithm to find permutations of rows and columns of $\diagup$-free matrix to avoid $\diagup$.

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